

PostgreSQL 15 and beyond

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# Evolution of the OSS database PostgreSQL



- Ongoing version upgrades once a year
- Enhanced support for large volume data in recent years





# Agenda

- Key features in PostgreSQL 15
- Performance improvements in PostgreSQL 15
- PostgreSQL 16 and beyond



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#### Key features in PostgreSQL 15

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- Command handles inserts, updates, and deletes, all in a single transaction
- Attempts to insert a new product if the given products already exists, then update those, otherwise delete them

```
MERGE INTO TargetProducts Target
USING SourceProducts Source
ON Source.ProductID = Target.ProductID
WHEN NOT MATCHED AND Source.ProductId IS NOT NULL THEN
INSERT VALUES (Source.ProductID, Source.ProductName, Source.Price)
WHEN MATCHED AND Target.ProductName IN ('Table', 'Desk') THEN
UPDATE SET ProductName = Source.ProductName, Price = Source.Price
WHEN MATCHED THEN
DELETE;
```





- Ensures that the join produces <u>at most</u> one candidate change row for each target row
  - Otherwise, it will lead to error "MERGE command cannot affect row a second time"
- For each candidate change row, the first clause to evaluate as true is executed
- No more than one WHEN clause is executed for any candidate change row
- One common use case is while trying to maintain Slowly Changing Dimensions (SCD) in a data warehouse. In such cases, one needs to:
  - Insert new records into the data warehouse,
  - Remove records from the warehouse which are not in the source anymore, and
  - Update values in the warehouse which have been updated in the source

#### **Base backups**



- Allows to specify targets for backups
  - Backup location can be specified using -t target or --target=target
    - client: Default value
    - server: Stores backup on server
    - **blackhole**: Discards the contents, used only for testing and debugging purpose
  - This option cannot be used with default WAL streaming option -Xstream
- Backups can be compressed
  - Server-side compression
  - Client-side compression
  - Compression options: gzip, LZ4, and Zstandard
  - Client-side: gzip was supported prior to 15
  - Allows faster and smaller backups





- Full Page Writes can be compressed using LZ4 and Zstandard compression
  - Previously, PGLZ was used
  - User can specify the compression method via GUC wal\_compression
- WAL archiving via loadable modules
  - Currently, the continuous archiving is done via shell commands
  - This new feature allows users to create custom modules for continuous archiving
  - Although archiving via a shell command is much simpler, a custom archive module will often be considerably more robust and performant
  - For example, if the archiving destination path already exists and has contents identical to source path, we can allow archiving to succeed via custom module
  - Users can use GUC archive\_library to specify a library that can be used to archive a logfile segment





- Stores statistics in shared memory
  - No more UDP transfer for stats
  - No more writing to temp files
  - The stats in shared memory will be written once at server shutdown
  - The prior stats state is restored after shutdown and restart of server
  - The stats are discarded on restarting server after a crash
- Allows log output in JSON format
  - log\_destination=jsonlog
  - This is the 3<sup>rd</sup> type of destination of this kind, after stderr and csvlog
  - The format is convenient to feed logs to other applications







- New stats in pg\_stat\_statements
  - I/O timing for temp files
  - JIT counters
- Extended statistics (CREATE STATISTICS ...) to record statistics for a parent with all its children
  - Regular statistics already tracked this information
  - Improve queries that involve processing the inheritance tree as a whole





• Adds support for prepared transactions to built-in logical replication

CREATE PUBLICATION mypub FOR ALL TABLES; CREATE SUBSCRIPTION mysub CONNECTION 'dbname=postgres' PUBLICATION mypub WITH (two phase = true);

- Reduces the lag to replicate data
- This provides the base to build conflict-free logical replication
- Allows publication of all tables in a schema

CREATE PUBLICATION mypub FOR TABLES IN SCHEMA mysch; CREATE PUBLICATION mypub FOR TABLE mytab, TABLES IN SCHEMA mysch;

• Tables added later to the listed schemas will also be replicated





• Allows publication content to be filtered using a WHERE clause

CREATE PUBLICATION mypub FOR TABLE mytab WHERE (c1 > 10);

- This can help distribute data among nodes and improve performance by sending data selectively
- The WHERE clause allows only simple expressions
  - It cannot contain user-defined functions, operators, types, collations, system column references or non-immutable built-in functions
- If a publication publishes UPDATE or DELETE operations, the row filter WHERE clause must contain only columns that are covered by the replica identity
- If a publication publishes only INSERT operations, the row filter WHERE clause can use any column





Allows publications to publish specific columns for tables

CREATE PUBLICATION mypub FOR TABLE mytab (c1, c2);

- The choice of columns can be based on behavioral or performance reasons
- A column list can contain only simple column references
- A column list can't be specified if the publication also publishes FOR TABLES IN SCHEMA
- If a publication publishes UPDATE or DELETE operations, any column list must include the table's replica identity columns
- If a publication publishes only INSERT operations, then the column list may omit replica identity columns
- Allows logical replication to run as the owner of the subscription
  - Only superusers, roles with bypassrls, and table owners can replicate into tables with row-level security policies





- Conflict resolution
  - Current methods:
    - By manually removing the conflicting data
    - By skipping the transaction via pg\_replication\_origin\_advance
  - A new method:
    - By specifying LSN of the conflicting transaction

ALTER SUBSCRIPTION mysub SKIP (lsn = 0/14C0378)

- Information related to LSN of failed transaction will be available in server logs
- Users can set a parameter to automatically disable replication on conflict
  - Useful for scenarios where a retry could not possibly succeed without human intervention
- New system view pg\_stat\_subscription\_stats
  - Shows stats about errors which occurred during the application of logical replication changes or during initial table synchronization





Allows ICU collations to be set as the default for clusters and databases

```
initdb --locale-provider=icu --icu-locale=en
CREATE DATABASE dbicu LOCALE_PROVIDER icu ICU_LOCALE 'en-u-kf-upper'
```

Allows unique constraints and indexes to treat NULL values as not distinct

```
CREATE TABLE mytab(c1 text UNIQUE NULLS NOT DISTINCT);
INSERT INTO mytab VALUES('amit');
INSERT INTO mytab VALUES(NULL);
INSERT INTO mytab VALUES(NULL);
ERROR: duplicate key value violates unique constraint "mytab_c1_key"
```

DETAIL: Key (c1)=(null) already exists.





• ON DELETE, partially SET NULL

• This is useful for multitenant or sharded schemas, where the tenant or shard ID is included in the primary key of all tables but shouldn't be set to null.





CREATE DATABASE new option – STRATEGY

#### WAL\_LOG

- Database will be copied block by block, and each block will be separately written to the WAL
- Checkpoints are not required
- Avoids the impact on overall system due to checkpoint
- Efficient strategy for small databases
- Default

#### • FILE\_COPY

- Current method
- Writes a small record to the write-ahead log for each tablespace used by the target database
- Each such record represents copying an entire directory to a new location at the filesystem level
- Reduces WAL volume for the large template database
- Performs checkpoint both before and after operation, which can impact system performance





- Granting SET and ALTER SYSTEM privileges for superuser GUC parameters
  - This allows to set superuser server parameters via non-superuser roles

GRANT SET ON PARAMETER wal\_compression TO bob; GRANT ALTER SYSTEM ON PARAMETER wal\_compression TO bob;

- Predefined role pg\_checkpoint
  - Allows members to run CHECKPOINT
  - Previously, checkpoints could only be run by superusers
- Security invoker views
  - Checks permissions for base relations using the privileges of the user of the view
  - Default is to check using the privileges of the view owner

CREATE VIEW myview WITH (security\_invoker=true) AS SELECT \* FROM mytab;



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- Improved performance and reduced memory consumption of in-memory sorts
  - Improved performance of single column sorts by more than 25%
    - Only when the result also contains single column
    - Will be used for SELECT col1 from mytab ORDER BY col1;
    - Will not be used for SELECT col1, col2 from mytab ORDER BY col1;
  - Reduced memory consumption by using generation memory context
    - We were using memory allocation scheme that rounded requests to the next power of 2
    - Performance improvement depends on tuple size, but up to ~40% improvement is observed
  - Reduced function call overhead by adding specialized sort routines for common datatypes
    - Performance improvement: ~5%





- Improved performance for sorts that exceed work\_mem
  - Switched to a batch sorting algorithm that uses more output streams than before
  - Performance improvement depends on work\_mem the smaller the work\_mem, the greater is the improvement
  - Performance improvement: ~40%
- For further details, see <u>blog</u>





- Allows parallel commit on postgres\_fdw servers
  - Enabled by the parallel\_commit option via CREATE/ALTER SERVER
  - Commits the transaction in parallel on all foreign servers involved in local transaction
  - This can improve performance of distributed PostgreSQL clusters using postgres\_fdw

#### Partitioning



- Ordered scans of partitions in more cases
  - Allows optimization when **DEFAULT** and **LIST** partitions containing multiple values gets pruned
  - Uses Append instead of MergeAppend
  - Improved performance by avoiding the need to sort
- Improved planning time for statements where only few of the partitions are relevant





- Parallelized SELECT DISTINCT
  - Introduces two-phase DISTINCT
  - Phase 1 is performed on parallel workers
    - Rows are made distinct there either by hashing or by sort/unique
  - Phase 2 is performed by leader backend
    - Removes duplicate rows that appear due to combining rows for each of the parallel workers

```
SELECT DISTINCT four FROM tenk1;

QUERY PLAN

Unique

-> Sort

Sort Key: four

-> Gather

Workers Planned: 2

-> HashAggregate

Group Key: four

-> Parallel Seq Scan on tenk1
```

#### Recovery/Replay



- Speeded up recovery/replay by prefetching needed file contents
  - recovery\_prefetch: When enabled, looks ahead in the WAL and try to initiate asynchronous reading of referenced data blocks not yet cached in our buffer pool
  - Works where posix\_fadvise() is available
  - wal\_decode\_buffer\_size: Maximum distance to read ahead in the WAL to prefetch referenced data blocks





- Allows vacuum to be more aggressive in setting the oldest frozenxid
  - Before 15, vacuum set it to whatever value was used to determine which tuples to freeze the FreezeLimit cutoff
  - Now, we set it to value <= the oldest extant XID remaining in the table</p>
  - This can be much more recent than FreezeLimit
  - This will help in reducing the times anti-wraparound vacuum is invoked for certain workloads

#### Changes in PostgreSQL 15



- New features and enhancements
  - Support for the SQL MERGE command
  - Selective publication of tables' contents within logical replication publications, through the ability to specify column lists and row filter conditions
  - More options for compression, including support for Zstandard (zstd) compression
    - Includes support for performing compression on server side during pg\_basebackup
  - Support for structured server log output using the JSON format
  - Performance improvements, particularly for in-memory and on-disk sorting
- The full list of new/enhanced features and other changes can be found <u>here</u>

# FUJITSU

## Agenda

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Disclaimer: This section is based on what I could see being proposed in community at this stage

#### PostgreSQL 16 and beyond





- Various improvements in Logical Replication
  - Allow same table replication by filtering based on origins
  - Parallel Apply
  - Replication of sequences
  - Enable logical replication from standby
  - DDL Replication
  - Use of indexes on subscriber when publisher has specified replica identity full
  - Replication of other objects like LOBs
- Reduced number of commands that need superuser privilege
- SQL/JSON improvements to make it more standard compliant

• ...

#### PostgreSQL 16 and beyond





- Transparent column encryption
  - Automatic, transparent encryption and decryption of particular columns in the client
- Change build infrastructure by replacing it with Meson build system
  - Developer-oriented feature
- Asynchronous I/O
  - Will allow prefetching data and will improve system performance
- Direct I/O
  - Will bypass the OS cache and lead to better performance in some cases
- Various improvements in Hash indexes
  - Allow unique indexes
  - Allow multi-column indexes

#### PostgreSQL 16 and beyond





- Improvements in vacuum technology by using performance data structure, advancing relfrozenxid earlier, and by reducing WAL volume
- Improvements in partitioning technology
- Improve statistics/monitoring
- 64bit XIDs
  - Can avoid freezing and reduce the need of autovacuum
- TDE
  - Can help in meeting security compliance in many organizations
- Incremental maintenance of materialized views

# Thank you

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